

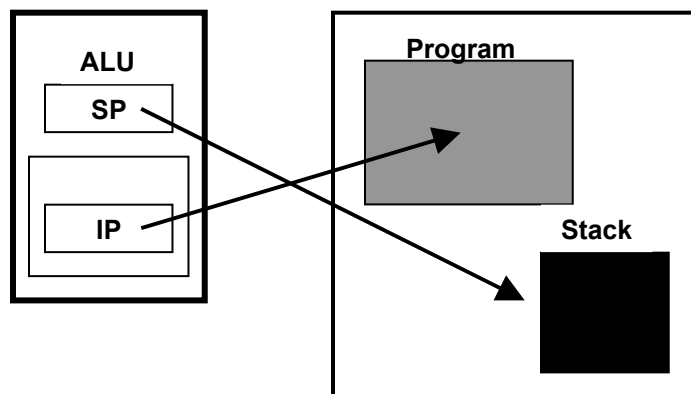
## CPU

The **CPU** is usually organised in three parts: The **Arithmetic and Logic Unit** or **ALU** (which performs computations and comparisons) and the **Control Unit** which fetches information from and stores information into the **Memory Unit**.

Every **CPU** has at least up to a dozen **registers** which are special, small, very fast memory.

**Registers** that can be used for any type of arithmetic, as well as for addresses (pointers, subscripts, etc...) are called **general-purpose Registers** e.g. **AX, BX, CX** and **DX**.

**Two registers** of invaluable importance are the **Stack Pointer (SP)** in the **ALU** and the **Instruction Pointer (IP)** also referred as the **Program Counter (PC)** in the **Control Unit**.



The **Instruction Pointer** contains the **address** in memory of the next instruction to be executed by the CPU. One of the main functions of the **Control Unit** is to perform an infinite loop of fetching the instruction at the **IP**, updating the **IP** to point to the next instruction and then executing the fetched instruction. This is called the **fetch-execute cycle**.

The **stack data structure** is a block of memory locations, and the **Stack Pointer** register in the **ALU** contains the address of one of them, called the '**top**' of the stack.

## IBM PC and compatible Memory Structure

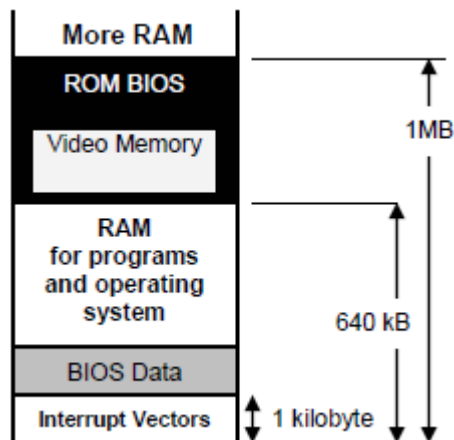
**Any IBM PC** makes special use of parts of its first **1 MB** address space.

**Any Intel 80X86 chip** uses the first **1024 bytes (1 kB)** for **interrupt vectors**.

The PC itself makes special use of memory in locations above the first **640 kB**.

**All IBM PCs** start up in **real mode** in which only the first **1 MB of RAM** is usable.

**All modern operating systems** immediately change to **protected mode** where **all the memory** becomes available.



## Memory Addressing

Two types of memory addressing: **16-bit** and **32-bit**

### ❖ 16-bit addressing:

The original IBM PC (80286) used this mode. However with 16-bit we can only reference  $2^{16}$  i.e. 65,536 bytes, 64KB of memory.

**So, how can we get full 20-bit address to be able to access the first 1MB of memory?**

Use segment + some offset ... *More later*

### ❖ 32-bit addressing

Used by the 80386 and later. This allowed us a potential  $2^{32}$  i.e. 4 GB of memory

Modern Pentium processors can address even up to 64 GB of memory. Therefore, 32-bit addressing seems like bliss as it breaks the 1 MB barrier but ...

**Issues:**

**Backwards compatibility:** 386 and later must still run programs requiring 16-bit addressing.

**Memory protection:** There must be a scheme to prevent programs from clobbering each other.

## Number Systems

### ❖ Denary (decimal) number system

This number system which is usually used by humans. Positional number system: the furthest right number has a weight of  $10^0$  e.g. decimal number 46,536 is

$$4 \times 10^4 + 6 \times 10^3 + 5 \times 10^2 + 3 \times 10^1 + 6 \times 10^0$$

### ❖ Binary number system

This number system which is usually used by computers. Positional number system: the furthest right number has a weight of  $2^0$  e.g. binary number 1011010B is

$$1 \times 2^6 + 0 \times 2^5 + 1 \times 2^4 + 1 \times 2^3 + 0 \times 2^2 + 1 \times 2^1 + 0 \times 2^0 = 64 + 16 + 8 + 2 = 90_{10}$$

**Conversions:**

Binary to denary: as above

Denary to binary: Reverse of Horner's method: repeated division by 2, sequence of remainders taken in reverse order.

### ❖ Hexadecimal number system

Binary becomes too cumbersome for e.g. 1 million in denary is 20 binary digits

Hence, hexadecimal (base 16) is used. Its digit starts from 0-9 and A-F.

**Conversions:**

Hex to Bin: each hex digit corresponds to exactly four binary digits.

Bin to Hex: Group of four bits taken from the right and padding with zero where needed.

Dec to Hex: divide repeatedly by 16, sequence of remainders taken in reverse order.

Hex to Dec: As denary, replacing base with 16.

## Assembler Overview

In this section will show you how to **write, assemble, link and execute** a program that displays your name on the screen.

As we will be using a **Command Prompt** window and hence using **virtual 8086** mode for all our programs, the naming convention for filenames must be of the **8.3 format**.

Anyway, even if MS Windows allow us **256-character** long filenames, the trademark of a good programmer is still 8.3 format filenames.

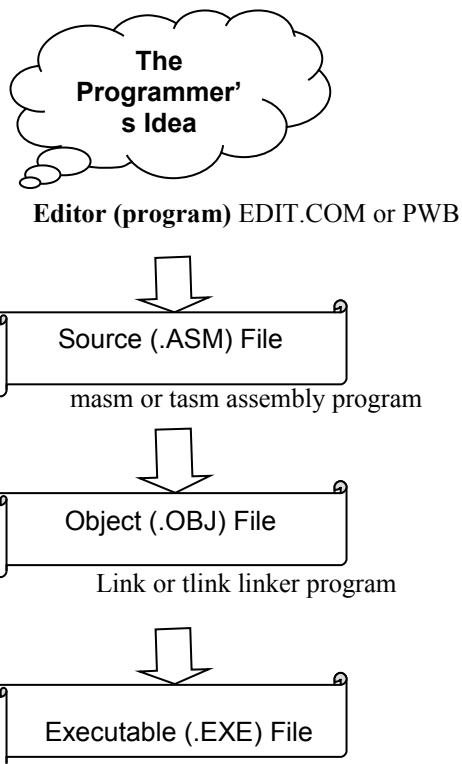
No blank space allowed in the filename. Use underscores (another trademark of an IT literate person!!!)

The **3 types of files** we will be working with are indicated by the following extensions:

- .ASM** Assembly language source programs that you write.
- .OBJ** Object files; the output produced by the assembler from an input **.ASM** file.
- .EXE** Executable files produced by the linker utility combining one or more **.OBJ** files.

Registers Discussed

EAX	ah AX
EBX	bh BX
ECX	ch CX
EDX	dh AX
ESI	SI
EDI	DI
EBP	BP
	SP
	DS
	ES
	FS
	GS
	SS
	CS
	IP
	o d i t s z



**:: FIRST.ASM – Our first Assembly Language Program. This program displays the line “ Hello, my name is nefertum” on the screen**

```

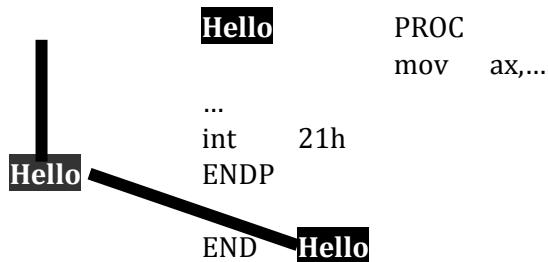
.MODELSMALL
.586           ; allows Pentium instructions. Must come after .MODEL

.STACK        100h
.DATA
Message      DB      'Hello, my name is nefertum', 13, 10, '$'

.CODE
Hello        PROC
mov          ax, @data
mov          ds, ax
mov          dx, OFFSET Message
mov          ah, 9h           ; Function code for Display String
int          21h             ; The standard way to call the OS
mov          al, 0           ; Return code for 0
mov          ah, 4ch         ; Function code for Exit to OS
int          21h
Hello        ENDP
END          Hello           ; tells where to start execution
    
```

**Synopsis**

- ❖ Upper & lower case are not distinguished in Assembly Language (except within quotes).
- ❖ The name chosen for the **PROC**edure, **Hello** in *first.asm* is unimportant, but the three occurrence of **Hello** must match exactly.



- ❖ The **PROC** name **Hello** is independent from the source filename *first.asm*
- ❖ Anything following a semicolon on a line is a **comment** and is ignored by the assembler.

**General structure of an IBM PC Assembly Language**

Label	OPERATION	OPERANDS	;comments
-------	-----------	----------	-----------

★ **LABEL**

**Labels** are **optional** and are example of *symbolic constants*, which are strings of characters which the programmer invents. They can be up to 31 characters and can consists of alphanumeric and special characters like '@', '\_', '\$' and '?' but must not start by a number (to distinguish symbols from numbers).

## ▶ OPERATION

**2 types:** Actual (executable) single machine instructions, such as *mov*  
Non-executable instructions to the assembler, such as *.DATA, DB*

**The second type of operation is called a *pseudo-operation*.**

Both operations are specified using a symbolic code called the **OPERATION CODE (OPCODE)**

*Pseudo-operations* are usually shown in **upper case** to distinguish them from executable machine instructions.

## ▶ OPERANDS

Each *op-code* can have zero or more operands. **Real** machine instructions usually have one or two.

## Global Program Structure

**:: FIRST.ASM – Our first Assembly Language Program. This program displays the line “ Hello, my name is nefertum” on the screen**

	.MODEL	SMALL	
	.586	<b>; allows Pentium instructions. Must come after .MODEL</b>	
	.STACK	100h	Stack Segment
Message	.DATA		Data Segment
	DB	'Hello, my name is nefertum', 13, 10, '\$'	
Hello	.CODE		Code Segment
	PROC		
	mov	ax, @data	
	mov	ds, ax	
	mov	dx, OFFSET Message	
	mov	ah, 9h	<b>; Function code for Display String</b>
	int	21h	<b>; The standard way to call the OS</b>
	mov	al, 0	<b>; Return code for 0</b>
	mov	ah, 4ch	<b>; Function code for Exit to OS</b>
	int	21h	
Hello	ENDP		
	END	Hello	<b>; tells where to start execution</b>

The **Stack Segment** is used to reserve space for the **stack**. All programs must use a **stack**.

The **Data Segment** is used to declare variables and data storage, with/without initialization.

The **Code Segment** is used for executable program code.

## Synopsis

### **.MODEL**

This defines which **memory model** to use. Will be discussed later. Assume **SMALL** for now.

### **.586**

This allows **Pentium instructions**; none are really used in this program. Will be discussed later.

## .STACK Segment

The .STACK segment is declared with the statement

```
.STACK      n
```

where **n (which must be even)** is the number of bytes reserved for the stack. We will use  $n = 100h$  i.e.  $256_{10}$  which should be enough for our purposes.

## .DATA Segment

The .DATA segment is used to define variables, with or without initial values. In our first, program we have used:

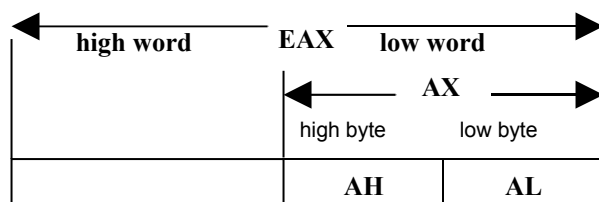
```
Message     DB    'Hello, my name is nefertum', 13, 10, '$'
```

uses **DB** (*define byte*) pseudo-operation to define a string of byte, which is the message to be displayed on the screen. Three commonly used data types are **byte**, **word** and **double word**. These are declared respectively by **DB**, **DW** and **DD**. These pseudo-operations also allow us to set initial values for the variables. E.g.

```
Unlucky     DB    13
Minus       DW   -299
Million     DD   1000000
```

defines *Unlucky* to be a byte (8-bit) variable with an initial value of 13, *Minus* to be a word (16-bit) variable with an initial value of -299 and *Million* to be a double word (32-bit) variable with an initial value of 1000000. Trying to initialize a value outside these ranges will cause an assembler error.

The **80X86 ALU** (Arithmetic and Logic Unit) has several registers, among which are four important registers used for computations. In the **286**, they were limited to **16 bits** and called **AX**, **BX**, **CX** and **DX** registers. From the **386** and later, there were extended to double words of **32 bits**, **EAX**, **EBX**, **ECX** and **EDX** respectively. The format of the **EAX** register is as follows:



The other registers are similar. “low” and “high” refer to the value they contain- the bits in the high byte represent larger value than those in the low byte (**Big-Endian**). There is no name for the high (left) word and there is no way to access it directly.

The **Stack Pointer SP** and the **Instruction Pointer IP** are **16 bit registers** that are **addresses in memory** and also have their extended (32-bit mode) counterparts' viz. **ESP** and **EIP**

The **executable** part (.CODE segment) of the **Hello** program does **three** things:

mov	ax, @data	<b>Standard code</b>
mov	ds, ax	

mov	dx, OFFSET Message	<b>DOS Display String Call</b>
mov	ah, 9h	
int	21h	

mov	al, 0	<b>DOS Return to DOS Call</b>
mov	ah, 4ch	
int	21h	

Two instructions are used: the **mov** instruction which moves data from place to place and the **int 21h** instruction which causes the operating system to perform various functions. Most 80X86 **opcodes** are written in three letters.

❖ The **mov** instruction

Syntax:

mov    *destination, source*                    ;destination = source

**mov** is really a copy instruction since the previous content of the source is unchanged. Thus if **ax** = 1234 and **bx** = 9876 and the instruction **mov ax, bx** is executed, we will have both **ax** and **bx** equal to 9876.

There are various possibilities for the source and destination operands:

**mov**    *memory or register, memory or register or constant*

which is abbreviated as:

mov    *mem / reg, mem / reg / constant*

**Note:** **mov** copies FROM right TO left-hand operand.

## GOLDEN RULES

1. You cannot move MEMORY to MEMORY directly
2. You cannot move TO a CONSTANT
3. You cannot move operands of DISSIMILAR SIZES

E.g.

mov	ax, A	;ok
mov	A, ax	;ok
mov	ax, bx	;ok
mov	A, 345	;ok
mov	ax, 345	;ok
mov	A, B	;illegal – memory to memory
mov	345, A	;illegal – can't move TO constant
mov	ax, 24	;ok
mov	al, 24	;ok
mov	eax, 24	;ok
mov	al, 345	;illegal – word to byte move
mov	ax, 100000	;illegal – double to word move
mov	ax, al	;illegal – byte to word move

❖ The **int 21h** instruction

Operating system functions on the IBM PC are called DOS calls, after the original operating system MS-DOS. Using DOS is much like using standard procedures like puts(), printf(), cout in C/C++.

**DOS Call Syntax:**

```

; Load parameters to DOS Call into registers
mov ah, function code
int 21h ;call DOS and returned values (if any) are in registers

```

Each DOS call has its own **function code number** (in Hex) and pattern of registers used for sending and receiving values.

❖ **The DOS Display String Call**

The DOS call with function code 9h is used to display a string of characters on the screen. The 9h function displays a string of ASCII characters whose address is in the **dx** register, *up to but not including the first '\$' sign encountered*. The form of the call is

```

mov dx, OFFSET Message
mov ah, 9h
int 21h

```

The first **mov** instruction is used to load the address of the ASCII string into **dx**. The line in the data region:

```

Message DB "Hello, my name is nefertum", 13, 10, "$"

```

is used to define the string to be displayed. The numbers **13** and **10** are the ASCII values for **Carriage Return** and **Line Feed**, respectively.

❖ **The Exit to DOS Call**

An explicit return to the OS is needed because the computer will continue trying to execute whatever garbage lies in memory **beyond the end of your program**. To end execution of your program, you must execute the return to DOS call, which is:

```

mov al, 0h ; return code
mov ah, 4ch ; Exit to DOS function code
int 21h

```

The two **mov** instructions could also have been written as a single instruction:

```

mov ax, 4c00h

```

**DOS** call **4ch** has one parameter, the **return code** in the **ah** register. The return code is traditionally zero for a normal return, nonzero for an error return. This call is analogous to the C/C++ **exit()** function and the return statement in the **main()** function that returns a value.

## Down Memory Lane

A **byte** can contain any unsigned number from **0** to **255** or any signed number from **-128** to **127** (signed numbers **-128** to **-1** corresponds to the unsigned numbers **128** to **255** in a byte)

You can specify a sequence of values as such:

```
Abyte DB 12, 99, 20
```

defines 3 consecutive variables with initial values 12, 99 and 20 respectively. *Abyte* is the name of the *first variable* whose value is 12. The byte containing 99 is an 'anonymous' variable which can be referred as [*Abyte+1*] and the variable containing value 20 as [*Abyte+2*]

These 2 code fragments are equivalent:

```

      Abyte DB 12
      DB 99
      DB 20
Or
      Abyte DB 12, 99
      DB 20

```

## Uninitialized Variables

```

ByteVar DB ? ;ByteVar is a byte with NO initial value
WordVar DW ? ;WordVar is a word with NO initial value

```

## ASCII Variables

Bytes containing the ASCII representation of characters can be defined by enclosing the characters in paired single (') or double (") quotes:

```

      Hi DB 'H', 'e', 'l', 'l', 'o'
is equal to
      Hi DB 'Hello'
or
      Hi DB "Hello"

```

## Non-printable Characters

For non-printable characters, such as Carriage Return or Line Feed, (whose ASCII values are 13 and 10 respectively) the only way to get them with a DB instruction is to include them numerically:

```
Message DB 'Hello, my name is nefertum', 13, 10, '$'
```

## Creating Arrays

Data definitions can be repeated by using the DUP operation. E.g.

**Ones            DB    100 DUP (1)**

defines 100 consecutive bytes containing the number 1 (the first is called '**Ones**')

**DB    10 DUP (2, 3, ?)**

defines 30 bytes which are consecutively 2, 3, ?, 2, 3, ?, ...

## DW

DW operates similarly except that it defines words of storage.

**Aword            DW    1234h**

We can represent the result of as occupying two consecutive bytes of memory as follows:

e.g. 

12		34
----	--	----